Formal Abstractions for Packet Scheduling

Mohan, Liu, Foster, Kappé, Kozen





Early goal: routing.



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But now we need control over scheduling.



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Basic tools work fine...





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Early goal: routing.

But now we need control over scheduling.

Basic tools work fine...







R traffic goes to either Purmerend or The Hague.



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R traffic goes to either Purmerend or The Hague.

Goal:

Interleave R and B; interleave P and T.





R traffic goes to either Purmerend or The Hague.





3



R traffic goes to either Purmerend or The Hague.





3



New plan!

4



New plan! Interleave small, medium, and large packets.



New plan! Interleave small, medium, and large packets.







New plan! Interleave small, medium, and large packets.









A human needs a

range of trees.



A human needs a

range of trees.







A human needs a

range of trees.





A human needs a *range* of trees.





A human needs a *range* of trees.

this work





this work

A human needs a *range* of trees.

this work





Aside: PIFO Trees Sivaraman et al. at SIGCOMM '16



Just an ordered collection.





Just an ordered collection. Two ways of interacting with the collection:







Just an ordered collection. Two ways of interacting with the collection:

push







Just an ordered collection. Two ways of interacting with the collection:

push







Just an ordered collection. Two ways of interacting with the collection:

push







Just an ordered collection. Two ways of interacting with the collection:

push







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push









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Just an ordered collection. Two ways of interacting with the collection:

push









Review: priority queue

Everything from before holds, but we have a little more control.






Everything from before holds, but we have a little more control.







Everything from before holds, but we have a little more control.







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Everything from before holds, but we have a little more control.







Everything from before holds, but we have a little more control.









Review: priority queue The priority need not be inherent to the item!



































Introducing: PIFO Just a PQ, with a ranking function, but with rank-ties broken in FIFO order.





10

Traffic incoming from Rotterdam and Beverwijk





Traffic incoming from Rotterdam and Beverwijk

Goal: interleave R and B

11





Traffic incoming from Rotterdam and Beverwijk Goal: interleave R and B A PIFO will suffice.





Traffic incoming from Rotterdam and Beverwijk Goal: interleave R and B <u>A PIFO will suffice.</u>





Traffic incoming from Rotterdam and Beverwijk Goal: interleave R and B <u>A PIFO will suffice.</u>

$B_n, ..., B_1, (R,B)^*$





Traffic incoming from Rotterdam and Beverwijk Goal: interleave R and B

A PIFO will suffice.







Traffic incoming from Rotterdam and Beverwijk Goal: interleave R and B A PIFO will suffice. \rightarrow B₁, (R,B)*







Traffic incoming from Rotterdam and Beverwijk Goal: interleave R and B <u>A PIFO will suffice.</u>

$B_n, ..., B_1, (R,B)^*$





Traffic incoming from Rotterdam and Beverwijk Goal: interleave R and B A PIFO will suffice.

Bn, ..., B1, (R,B)* Rn, ..., R1, (R,B)*







Traffic incoming from Rotterdam and Beverwijk Goal: interleave R and B A PIFO will suffice.

Bn, ..., B1, (R,B)* Rn, ..., R1, (R,B)* (R,B)*





R traffic goes to either Purmerend or The Hague.



R traffic goes to either Purmerend or The Hague. Goal:

Interleave R and B; interleave P and T.





B3, B2, P2, B1, P1



$T_1 \longrightarrow B_3, B_2, P_2, B_1, P_1$



$T_1 \longrightarrow B_3, B_2, P_2, B_1, P_1$

B₃, B₂, P₂, T₁, B₁, P₁



$T_1 \rightarrow B_3, B_2, P_2, B_1, P_1$

B₃, B₂, P₂, T₁, B₁, P₁ B₃, B₂, P₂, B₁, T₁, P₁



$T_1 \rightarrow B_3, B_2, P_2, B_1, P_1$

B₃, B₂, P₂, T₁, B₁, P₁ B₃, B₂, P₂, B₁, T₁, P₁ B₃, B₂, P₂, B₁, P₁, T₁



$T_1 \longrightarrow B_3, B_2, P_2, B_1, P_1$ B3, B2, $B_{1},$ B3, B2, B3, B2, , **Б**1,



$T_1 \longrightarrow B_3, B_2, P_2, B_1, P_1$

 $B_3, B_2, P_2, T_1, B_1, P_1$ B₃, B₂, P₂, B₁, T₁, P₁ B₃, B₂, P₂, B₁, P₁, T₁

14
$T_1 \longrightarrow B_3, B_2, P_2, B_1, P_1$

B3, B2, P2, T1, B1, P1 B₃, B₂, P₂, B₁, T₁, P₁ B3. P2, P2, B1, P1, T1



B_3, B_2, P_2, B_1, P_1

B₃, B₂, P₂, T₁, B₁, P₁ B₃, B₂, P₂, B₁, T₁, P₁ B₃, P₂, P₂, B₁, P₁, T₁

 $B_3, P_2, B_2, T_1, B_1, P_1$



$F_1 \longrightarrow B_3, B_2, P_2, B_1, P_1$

B3, B2, P2, T1, B1, P1 B₃, B₂, P₂, B₁, T₁, P₁ B₃, P₂, P₂, B₁, P₁, T₁

 $B_3, P_2, B_2, T_1, B_1, P_1$ B3, P2, B2, P1, B1, T1



$T_1 \longrightarrow B_3, B_2, P_2, B_1, P_1$



B₃, B₂, P₂, T₁, B₁, P₁ B3, B2, P2, B1, T1, P1 B₃, P₂, P₂, B₁, P₁, T₁



$F_1 \longrightarrow B_3, B_2, P_2, B_1, P_1$

B3, B2, P2, T1, B1, P1 B₃, B₂, P₂, B₁, T₁, P₁ B₃, P₂, P₂, B₁, P₁, T₁

 $B_3, P_2, B_2, T_1, B_1, P_1$ B3, P2, B2, P1, B1, T1



$F_1 \rightarrow B_3, B_2, P_2, B_1, P_1$

B₃, B₂, P₂, T₁, B₁, P₁ B₃, B₂, P₂, B₁, T₁, P₁ B₃, P₂, P₂, B₁, P₁, T₁

B₃, P₂, B₂, T₁, B₁, P₁ B₃, P₂, B₂, P₁, B₁, T₁



$T_1 \longrightarrow B_3, B_2, P_2, B_1, P_1$

B₃, B₂, P₂, T₁, B₁, P₁ B₃, B₂, P₂, B₁, T₁, P₁ B₃, P₂, P₂, B₁, P₁, T₁





Enqueueing a packet can require the reordering of buffered packets.

No PIFO can do this.









This behaves like a queue!





































B1, P1













This behaves like a queue! How do we pop it? How do we push into it?











interleave R and B; interleave P and L

push T₁







interleave R and B; interleave P and L

push T₁







interleave R and B; interleave P and L

push 1







interleave R and B; interleave P and L

push 1

























R traffic goes to either Purmerend or The Hague. Goal:

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Key Insight

A PIFO tree manifests a programming language.




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A PIFO tree manifests a programming language.





A PIFO tree manifests a programming language.

A program is precisely a scheduling algorithm.

language tree expressivity shape





tree shape

language expressivity



tree shape

Compare expressivity of languages?

language expressivity



tree shape

Compare expressivity of languages? Compare expressivity of trees?

language expressivity



tree shape

Compare expressivity of languages? Compare expressivity of trees?

Compile a program so it runs against a new tree?

language expressivity



A human needs a *range* of trees.



A human needs a **C** *range* of trees.



A human needs a **C** *range* of trees.

some sufficiently expressive tree





A human needs a **C** *range* of trees.

compilation

some sufficiently expressive tree





Formal model of PIFO trees



Formal model of PIFO trees w.r.t. tree shape

General theorems of expressiveness



Formal model of PIFO trees w.r.t. tree shape

Compiler

General theorems of expressiveness



Formal model of PIFO trees w.r.t. tree shape Compiler Simulator

General theorems of expressiveness



Trees with more leaves are more expressive. Taller trees are more expressive.



Trees with more leaves are more expressive. Taller trees are more expressive.

Captured elegantly by:



Trees with more leaves are more expressive. Taller trees are more expressive.

Captured elegantly by: Homomorphic embedding.

Map root to root, leaves to leaves. Respect ancestry.





Homomorphic embedding. Map root to root, leaves to leaves. Respect ancestry.





Homomorphic embedding. Map root to root, leaves to leaves. Respect ancestry.





Homomorphic embedding. Map root to root, leaves to leaves. Respect ancestry.


































































































Path: $[(2, r_1), ...]$ 1, 2, 3, 1, 2, 3, 1, 2, 3





Path: $[(2, r_1), ...]$ 1, 2, 3, 1, 2, 3, 1, 2, 3

























Homomorphic embedding. Map root to root, leaves to leaves. Respect ancestry.

Homomorphic embedding.

Two new algorithms, both starting with heterogeneous source trees.

Map root to root, leaves to leaves. Respect ancestry.

Homomorphic embedding.

Two new algorithms, both starting with heterogeneous source trees.

1. If target tree is regular d-ary for some d.

Map root to root, leaves to leaves. Respect ancestry.

Homomorphic embedding.

Two new algorithms, both starting with heterogeneous source trees.

1. If target tree is regular d-ary for some d. 2. If target tree is itself heterogeneous.

Map root to root, leaves to leaves. Respect ancestry.

But the hardware supports a regular-branching binary tree.

But the hard a regular-bra No problem. Here's how I

- No problem. Here's how I'll use that tree.
- But the hardware supports a regular-branching binary tree.

No problem.

Simulation

$$PUSH(p, pkt, r) = p'$$

$$push(Leaf(p), pkt, r) = Leaf(p')$$

$$push(Interprete for the second se$$

 $push(qs[i], pkt, pt) = q' \qquad push(p, i, r) = p'$ nternal(qs, p), pkt, (i, r) :: pt) = lnternal(qs[i/q'], p')

$$PUSH(p, pkt, r) = p'$$

$$push(Leaf(p), pkt, r) = Leaf(p')$$

$$push(Interprete for the second se$$

 $push(qs[i], pkt, pt) = q' \qquad push(p, i, r) = p'$ iternal(qs, p), pkt, (i, r) :: pt) = Internal(qs[i/q'], p') Path

$$PUSH(p, pkt, r) = p' \qquad p$$

$$push(Leaf(p), pkt, r) = Leaf(p') \qquad push(Integer p)$$

 $\frac{ts \in \text{Topo}^n}{(i,r) :: pt \in \text{Path}(\text{Node}(ts))} \quad t \in \text{Path}(ts[i])$

 $push(qs[i], pkt, pt) = q' \qquad push(p, i, r) = p'$ nternal(qs, p), pkt, (i, r) :: pt) = Internal(qs[i/q'], p')Path

$$PUSH(p, pkt, r) = p'$$

$$push(Leaf(p), pkt, r) = Leaf(p')$$

$$PIFOTree$$

$$push(Integers)$$

 $\frac{ts \in \text{Topo}^n}{(i,r) :: pt \in \text{Path}(\text{Node}(ts))} \quad t \in \text{Path}(ts[i])$

ush(qs[i], pkt, pt) = q'PUSH(p, i, r) = p'ternal(qs, p), pkt, (i, r) :: pt) = Internal(qs[i/q'], p')PathPIFOTree

$p \in PIFO(Pkt)$ $Leaf(p) \in PIFOTree(*)$

 $r \in \mathsf{Rk}$ $r \in Path(*)$

$$PUSH(p, pkt, r) = p'$$

$$push(Leaf(p), pkt, r) = Leaf(p')$$

$$PIFOTree$$

$$push(Interval)$$

 $n \in \mathbb{N}$ $ts \in \text{Topo}^n$ $p \in \text{PIFO}(\{1, \dots, n\})$ $\forall 1 \leq i \leq n. qs[i] \in \mathsf{PIFOTree}(ts[i])$ $Internal(qs, p) \in PIFOTree(Node(ts))$

 $ts \in \text{Topo}^n$ $1 \le i \le n$ $r \in \text{Rk}$ $pt \in \text{Path}(ts[i])$ $(i, r) :: pt \in Path(Node(ts))$

> ush(qs[i], pkt, pt) = q' PUSH(p, i, r) = p'ernal(qs, p), pkt, (i, r) :: pt) = Internal(qs[i/q'], p')Path PIFOTree

$\mathbb{N} ts \in \text{Topo}^n p \in \text{PIFO}(\{1, \dots, n\})$ $\forall 1 \le i \le n. \ qs[i] \in \text{PIFOTree}(ts[i])$	
$Internal(qs, p) \in PIFOTree(Node(ts))$	
	Торо
$1 \le i \le n$ $r \in \mathbf{Rk}$ p	$t \in Path(ts[i])$
$(i, r) :: pt \in Path(Node(ts))$ Topo	
$\operatorname{oush}(qs[i], pkt, pt) = q'$	PUSH(p, i, r) = p'
ternal(qs, p), pkt, (i, r) :: pt) = Internal(qs[i/q'], p')	
Path	PIFOTree

$$\frac{n \in \mathbb{N} \quad ts \in \mathsf{Topo}^{n}}{\mathsf{Node}(ts) \in \mathsf{Topo}}$$

$$\mathbb{N} \quad ts \in \mathsf{Topo}^{n} \quad p \in \mathsf{PIFO}(\{1, \dots, n\})$$

$$\forall 1 \leq i \leq n. \ qs[i] \in \mathsf{PIFOTree}(ts[i])$$

$$\mathsf{Internal}(qs, p) \in \mathsf{PIFOTree}(\mathsf{Node}(ts))$$

$$\mathsf{Topo}$$

$$1 \leq i \leq n \quad r \in \mathsf{Rk} \quad pt \in \mathsf{Path}(ts[i])$$

$$(i, r) :: pt \in \mathsf{Path}(\mathsf{Node}(ts))$$

$$\mathsf{Topo}$$

$$\mathsf{push}(qs[i], pkt, pt) = q' \quad \mathsf{PUSH}(p, i, r) = p'$$

$$\mathsf{ternal}(qs, p), pkt, (i, r) :: pt) = \mathsf{Internal}(qs[i/q'], p')$$

$$\mathsf{Path} \quad \mathsf{PIFOTree}$$

A general way to deploy PIFO trees





Let the hardware support some tree.





Let the hardware support some tree.



Let the human program against some tree.

Let the hardware support some tree.





Let the human program against some tree.

Let the hardware support some tree.





Compilable?

Let the human program against some tree.

Let the hardware support some tree.





Compilable?

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Let the hardware support some tree.





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Formal Abstractions for Packet Scheduling

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